Lehrstuhl für Informatik I
Effiziente Algorithmen und
wissensbasierte Systeme
Universität Würzburg

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Prof. Dr. Alexander Wolff
Dipl.-Inform. Philipp Kindermann

Homework Assignment #2
Computational Geometry (Winter Term 2014/15)

Exercise 1

Set R be a set of n axis-parallel rectangles in the plane. Two rectangles *intersect* if there is a point that is contained in both rectangles (including the boundary).

a) A special case of an intersection is an *overlap*, that is, an intersection in which no rectangle completely contains the other one. For the sake of simplicity, we first consider overlaps as in figure 1, that is, the top-right corner of the lower rectangle lies to the right of the top-right corner of the upper rectangle.

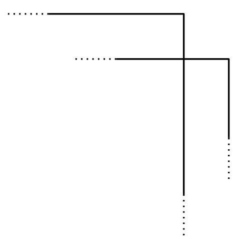

FIGURE 1: A special type of overlap.

Give an algorithm that decides in $O(n \log n)$ time *if* R contains two rectangles that overlap in the manner outlined. **[4 points]**

b) Give an algorithm that decides in $O(n \log n)$ time if R contains two intersecting rectangles. **[5 points]**

Hint: It is only important that the algorithm decides if there is *any* intersection. You do not have to determine the type of the intersection.

Exercise 2

Let P be a finite set of points in the plane. The *largest top-right region* of a point $p \in P$ is the union of all open axis-parallel square that touch p with their bottom-left corner and contain no point of P.

a) Prove that the largest top-right region of a point is either a square or the intersection of two open half-planes. **[2 points]**

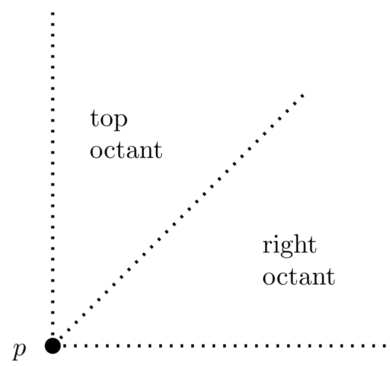

FIGURE 2: Top and right octant of the point p.

b) For every point $p \in P$ and each of the two corresponding octants (see Figure 2), consider which points of P restrict the largest top-right region of p the most. Is the knowledge of these points in both octants enough to determine the largest top-right region of p? **[2 points]**

c) Given a set P of n points, compute the largest top-right region of every point in p with total running time $O(n \log n)$. **[6 points]**

Hint: „Sweep" the plane twice to determine the points in subexercise b).

This assignment is due at the beginning of the next lecture, that is, on October 22 at 10:15. Solutions will be discussed in the tutorial on Friday, October 24, 14:00–15:30 in room SE I.